

A high-level implementation of software pipelining in LLVM

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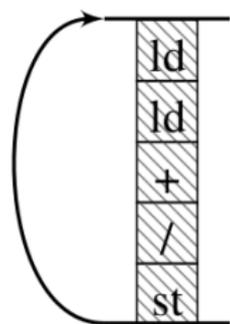
Rationale

Software pipelining (often Modulo Scheduling)

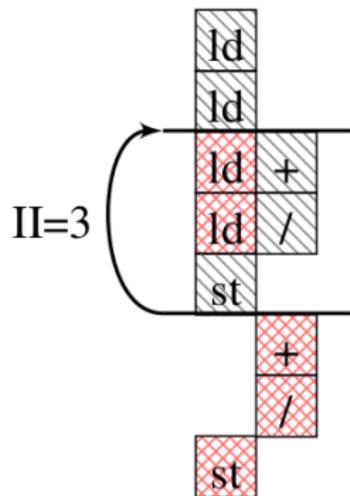
- ▶ Interleave operations from multiple loop iterations
- ▶ Improved loop ILP
- ▶ Currently missing from LLVM
- ▶ Loop scheduling technique
 - ▶ Requires both loop dependency and resource availability information
 - ▶ Usually done at a target specific level as part of scheduling
- ▶ But it would be very good if we could re-use this implementation for different targets

Example: resource constrained

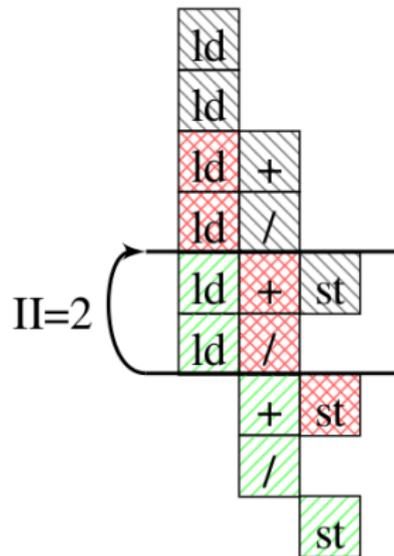
```
for(int i = 0; i < N; i++) {  
    B[i] = (A[2*i] + A[2*i+1]) / 2;  
}
```



(a) Original



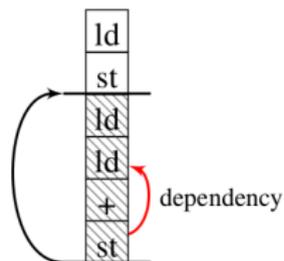
(b) Single memory



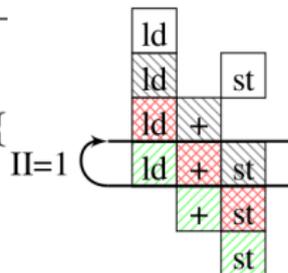
(c) Two memories

Example: data dependencies

```
B[0] = A[0];  
for(int i = 1; i < N; i++) {  
    B[i] = B[i-1] + A[i];  
}
```



```
register int r = A[0];  
B[0] = r;  
for(int i = 1; i < N; i++) {  
    r = r + A[i];  
    B[i] = r;  
}
```



Source Level Modulo Scheduling (SLMS)

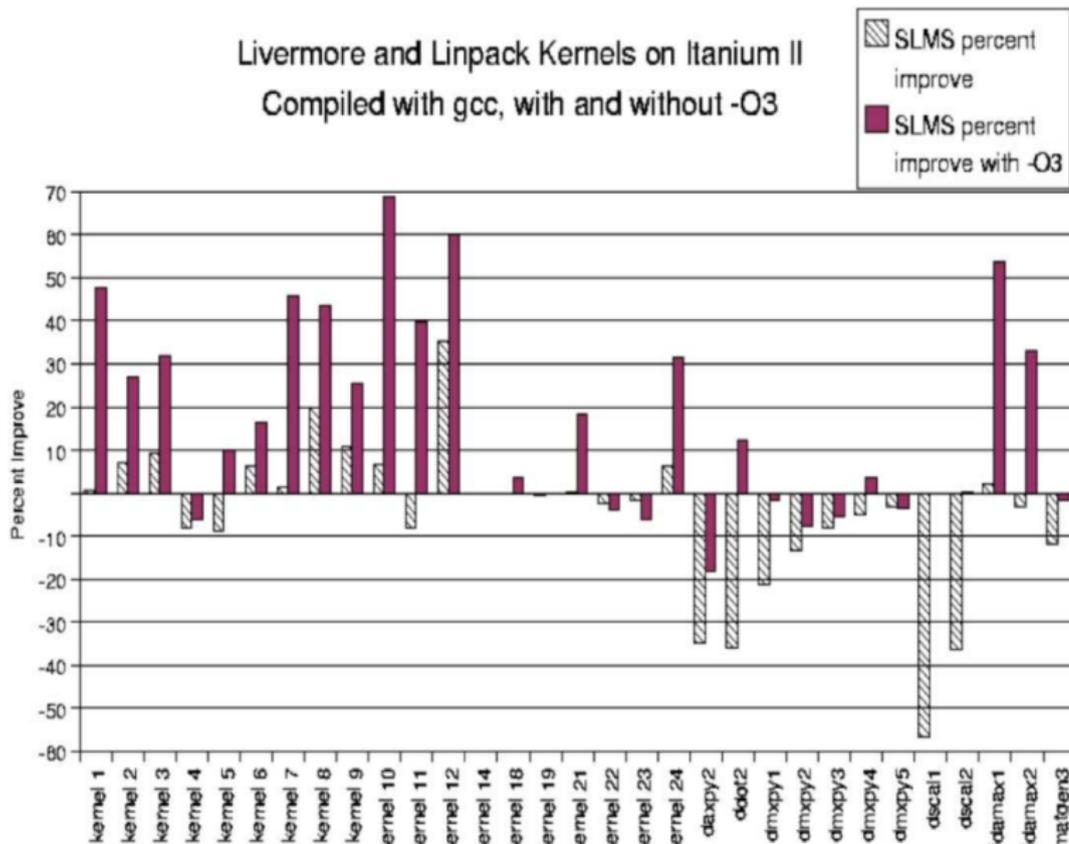
SLMS: Source-to-source translation at statement level

		$S1_0 : t = A[0] * B[0];$
$for(i = 0; i < n; i++)$		$for(i = 0; i < n - 1; i++)$
{		{
$S1_i : t = A[i] * B[i];$	\longrightarrow	$S2_i : s = s + t;$
$S2_i : s = s + t;$		$S1_{i+1} : t = A[i + 1] * B[i + 1];$
}		}
		$S2_{n-1} : s = s + t;$

Towards a Source Level Compiler: Source Level Modulo Scheduling
– Ben-Asher & Meisler (2007)

SLMS results

Livermore and Linpack Kernels on Itanium II
Compiled with gcc, with and without -O3



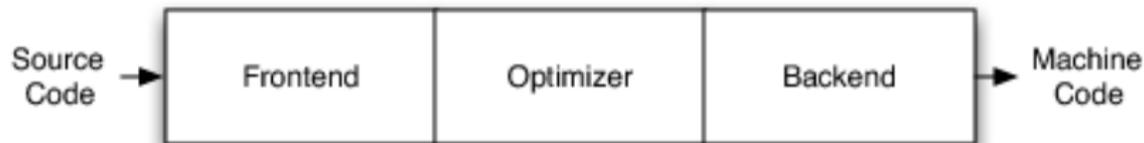
SLMS features and limitations

- ▶ Improves performance in many cases
- ▶ No resource constraints considered
- ▶ Works with complete statements
- ▶ When no valid II is found statements may be split (decomposed)

This work

What would happen if we do this at LLVM's IR level

- ▶ More fine grained statements (close to operations)
- ▶ Coarse resource constraints through target hooks
- ▶ Schedule loop pipelining pass late in the optimization sequence (just before final cleanup)



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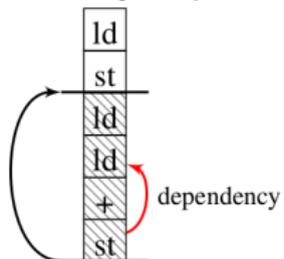
Implementation

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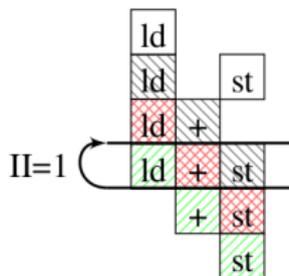
Conclusion

IR data dependencies

- ▶ Memory dependencies



- ▶ Phi nodes



Revisiting our example: memory dependencies

```
define void @foo(i8* nocapture %in, i32 %width) #0 {
entry:
    %cmp = icmp ugt i32 %width, 1
    br i1 %cmp, label %for.body, label %for.end

for.body:    ; preds = %entry, %for.body
    %i.012 = phi i32 [ %inc, %for.body ], [ 1, %entry ]
    %sub = add i32 %i.012, -1
    %arrayidx = getelementptr inbounds i8* %in, i32 %sub
    %0 = load i8* %arrayidx, align 1, !tbaa !0
    %arrayidx1 = getelementptr inbounds i8* %in, i32 %i.012
    %1 = load i8* %arrayidx1, align 1, !tbaa !0
    %add = add i8 %1, %0
    store i8 %add, i8* %arrayidx1, align 1, !tbaa !0
    %inc = add i32 %i.012, 1
    %exitcond = icmp eq i32 %inc, %width
    br i1 %exitcond, label %for.end, label %for.body

for.end:    ; preds = %for.body, %entry
    ret void
}
```

Revisiting our example: using a phi-node

```
define void @foo(i8* nocapture %in, i32 %width) #0 {
entry:
    %arrayidx = getelementptr inbounds i8* %in, i32 0
    %prefetch = load i8* %arrayidx, align 1, !tbaa !0
    %cmp = icmp ugt i32 %width, 1
    br i1 %cmp, label %for.body, label %for.end

for.body:    ; preds = %entry, %for.body
    %i.012 = phi i32 [ %inc, %for.body ], [ 1, %entry ]
    %0 = phi i32 [ %add, %for.body ], [ %prefetch, %entry ]
    %arrayidx1 = getelementptr inbounds i8* %in, i32 %i.012
    %1 = load i8* %arrayidx1, align 1, !tbaa !0
    %add = add i8 %1, %0
    store i8 %add, i8* %arrayidx1, align 1, !tbaa !0
    %inc = add i32 %i.012, 1
    %exitcond = icmp eq i32 %inc, %width
    br i1 %exitcond, label %for.end, label %for.body

for.end:    ; preds = %for.body, %entry
    ret void
}
```

Target hooks

- ▶ Communicate available resources from target specific layer
- ▶ Candidate resource constraints
 - ▶ Number of scalar function units
 - ▶ Number of vector function units
 - ▶ ...
- ▶ IR instruction cost
 - ▶ Obtained from CostModelAnalysis
 - ▶ Currently only a debug pass and re-implemented by each user (e.g. vectorization)

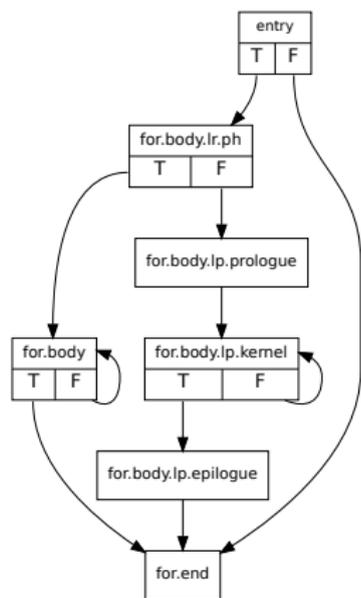
The scheduling algorithm

- ▶ Swing Modulo Scheduling
 - ▶ Fast heuristic algorithm
 - ▶ Also used by GCC (and in the past LLVM)
- ▶ Scheduling in five steps
 - ▶ Find cyclic (loop carried) dependencies and their length
 - ▶ Find resource pressure
 - ▶ Compute minimal initiation interval (II)
 - ▶ Order nodes according to 'criticality'
 - ▶ Schedule nodes in order

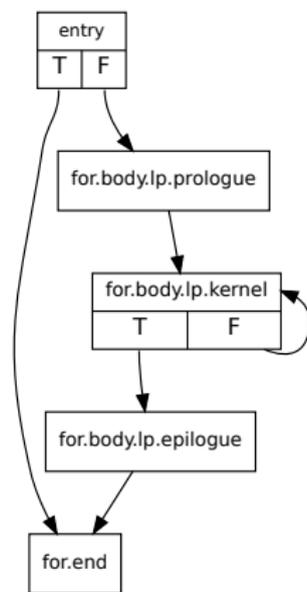
Swing Modulo Scheduling: A Lifetime-Sensitive Approach

– Llosa et al. (1996)

Code generation



CFG for 'loop5b' function



CFG for 'loop10' function

- ▶ Construct new loop structure (prologue, kernel, epilogue)
- ▶ Branch into new loop when sufficient iterations are available
- ▶ Clean-up through constant propagation, CSE, and CFG simplification

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Target platform

- ▶ Initial implementation for Movidius' SHAVE architecture
- ▶ 8 issue VLIW processor
- ▶ With DSP and SIMD extensions
- ▶ More on this architecture later today! (LG02 @ 14:40)
- ▶ But implemented in the IR layer so mostly target independent

Results

- ▶ Good points:
 - ▶ It works
 - ▶ Up to 1.5x speedup observed in TSVC tests
 - ▶ Even higher ILP improvements
- ▶ Weak spots
 - ▶ Still many big regressions (up to 4x slowdown)
 - ▶ Some serious problems still need to be fixed
 - ▶ Instruction patterns are split over multiple loop iterations
 - ▶ My bookkeeping of live variables needs improvement
 - ▶ Currently blocking some of the more viable candidate loops

Possible improvements

- ▶ User control
 - ▶ Selective application to loops (e.g. through #pragma)
- ▶ Predictability
 - ▶ Modeling of instruction patterns in IR
 - ▶ Improved resource model
 - ▶ Better profitability analysis
 - ▶ Superblock instruction selection to find complex operations crossing BB bounds?

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Conclusion

- ▶ It works, somewhat. . .
- ▶ IR instruction patterns are difficult to keep intact
- ▶ Still lots of room for improvement
 - ▶ Upgrade from LLVM 3.5 to trunk
 - ▶ Fix bugs (bookkeeping of live values, . . .)
 - ▶ Re-check performance!
 - ▶ Fix regressions
 - ▶ Test with other targets!

Thank you

TU/e

HIPEAC

Movidius 